

1 Abbreviation of complex types

unabbreviated	abbreviated	domain name
$\langle e, t \rangle$	et	D_{et}
$\langle e, e \rangle$	ee	D_{ee}
$\langle e, \langle e, t \rangle \rangle$	$\langle e, et \rangle$	$D_{e,et}$
$\langle \langle e, t \rangle, t \rangle$	$\langle et, t \rangle$	$D_{et,t}$
$\langle \langle e, t \rangle, e \rangle$	$\langle et, e \rangle$	$D_{et,e}$
$\langle d, \langle e, t \rangle \rangle$	$\langle d, et \rangle$	$D_{d,et}$
$\langle \langle e, t \rangle, \langle e, t \rangle \rangle$	$\langle et, et \rangle$	$D_{et,et}$
$\langle \langle e, t \rangle, \langle \langle e, t \rangle, t \rangle \rangle$	$\langle et, \langle et, t \rangle \rangle$	$D_{et,(et,t)}$

2 Abbreviation of Lambda-expressions

fully written	subscript type	no type
$\lambda x \in D_e. \text{dog}(x)$	$\lambda x_e. \text{dog}(x)$	$\lambda x. \text{dog}(x)$
$\lambda y \in D_e. \text{walk}(y)$	$\lambda y_e. \text{walk}(y)$	$\lambda y. \text{walk}(y)$
$\lambda z \in D_e. \text{orange}(z)$	$\lambda z_e. \text{orange}(z)$	$\lambda z. \text{orange}(z)$
$\lambda x \in D_e. \lambda y \in D_e. \text{take}(x)(y)$	$\lambda x_e \lambda y_e. \text{take}(x)(y)$	$\lambda x \lambda y. \text{take}(x)(y)$
$\lambda x \in D_e. \lambda z \in D_e. \text{see}(x)(z)$	$\lambda x_e \lambda z_e. \text{see}(x)(z)$	$\lambda x \lambda y. \text{see}(x)(z)$
$\lambda a \in D_e. \lambda b \in D_e. \text{hire}(a)(b)$	$\lambda a_e \lambda b_e. \text{hire}(a)(b)$	$\lambda a \lambda b. \text{hire}(a)(b)$
$\lambda f \in D_{\langle e, t \rangle}. \lambda y \in D_e. f(y)$	$\lambda f_{\langle e, t \rangle} \lambda y_e. f(y)$	$\lambda f \lambda y. f(y)$
$\lambda f \in D_{\langle e, t \rangle}. \lambda g \in D_{\langle e, t \rangle}. \exists x [f(x) = 1 \ \& \ g(x) = 1]$	$\lambda f_{\langle e, t \rangle} \lambda g_{\langle e, t \rangle}. \exists x [f(x) = 1 \ \& \ g(x) = 1]$	$\lambda f \lambda g. \exists x [f(x) = 1 \ \& \ g(x) = 1]$

3 Removal of = 1

(Abbreviate types and lambdas, while you're at it)

[[<i>the</i>]]	$\lambda f \in D_{\langle e, t \rangle}. \iota x \in C [f(x) = 1]$	$\lambda f_{et}. \iota x [f(x)]$
[[<i>red car</i>]]	$\lambda x \in D_e. \text{red}(x) = 1 \ \& \ \text{car}(x) = 1$	$\lambda x_e. \text{red}(x) \ \& \ \text{car}(x)$
[[<i>happy dog</i>]]	$\lambda x \in D_e. \text{happy}(x) = 1 \ \& \ \text{dog}(x) = 1$	$\lambda x_e. \text{happy}(x) \ \& \ \text{dog}(x)$
[[<i>every</i>]]	$\lambda f \in D_{\langle e, t \rangle}. \lambda g \in D_{\langle e, t \rangle}. \forall x [f(x) = 1 \rightarrow g(x) = 1]$	$\lambda f_{et} \lambda g_{et}. \forall x [f(x) \rightarrow g(x)]$
[[<i>no</i>]]	$\lambda f \in D_{\langle e, t \rangle}. \lambda g \in D_{\langle e, t \rangle}. \neg \exists z [f(z) = 1 \ \& \ g(z) = 1]$	$\lambda f_{et} \lambda g_{et}. \neg \exists z [f(z) \ \& \ g(z)]$
[[<i>some</i>]]	$\lambda f \in D_{\langle e, t \rangle}. \lambda g \in D_{\langle e, t \rangle}. \exists x [f(x) = 1 \ \& \ g(x) = 1]$	$\lambda f_{et} \lambda g_{et}. \exists x [f(x) \ \& \ g(x)]$